Simple Mechanisms for Welfare Maximization in Rich Ad Auctions

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Motivation – Rich Ad Auctions

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Packing Problem:

Ads have varying sizes

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- Multiple relevant ad formats per advertiser
- Platform has a size constraint

Model – Rich Ads

Publicly known \checkmark Size w_j for all ads $j \in S$ Expected no. clicks α_j for all ads $j \in S$

Set of rich ads A_i Value per click v_i



Unit-demand advertiser

Utility:

When $j \in A_i$ is allocated at price pthe utility of i is $v_i \alpha_j - p$

Model – Rich Ads

Publicly known \checkmark Size w_j for all ads $j \in S$ Expected no. clicks α_j for all ads $j \in S$

Reported rich ads $S_i \subseteq A_i$ Bid per click b_i

 (S_i, b_i)

Unit-demand advertiser

Utility:

When $j \in A_i$ is allocated at price pthe utility of i is $v_i \alpha_j - p$

Model – Rich Ad Auction

Feasibility Constraints:

- [UNIT-DEMAND] Allocate (at most) one ad $j(i) \in S_i$ per advertiser i
- [KNAPSACK] Total space allocated is $\sum_{i} w_{j(i)} \leq W$

Objective:

• Choose a feasible allocation to maximize SOCIAL WELFARE $\sum_{i} v_i \alpha_{i(i)}$

Algorithm Design

Simpler Approximation Algorithm

Optimal is NP-Hard!

[Sinha-Zoltners'79]

Best of *Incremental-Bang-per-Buck* allocation and highest value ad gives a 2-approximation to the optimal social welfare.

[This Paper]

Lemma. Best of *Bang-per-Buck* allocation and highest value ad gives a **3-approximation** to the optimal welfare.

--- Bang-per-Buck allocation: Allocate ads greedily in order of $\frac{b_i \alpha_j}{w_i}$ ----

Mechanism Design

Simple & Truthful Auction!

[This paper]

Main Theorem. There is a simple truthful mechanism which gets a 3-approximation to the optimal social welfare.

Monotone allocation rule

- Bang-per-Buck Allocation with prob. 2/3 and highest value ad with prob. 1/3
- Allocation (expected clicks) is monotone in both bid-per-click b_i and set of rich ads S_i



Myerson's lemma like

= Truthful Auction

Generalized Second Price Auction

- Our monotone allocation can be easily paired with GSP prices!
- GSP: The minimum bid b_i^* that won't change *i*'s allocation
- Very simple payments
- GSP is very widely used in practice in ad auctions
- GSP is not truthful!

Q. Welfare guarantee at Equilibrium?

Price of Anarchy of GSP

[This paper]

Theorem. Randomizing between the *Bang-per-Buck* allocation and highest value ad has a pure PoA of 6 and a Bayes-Nash PoA of $\frac{6}{\left(1-\frac{1}{e}\right)}$

Empirical Evaluation

Experimental Results

- On average our mechanism obtained 98% of the optimal welfare
- On average our mechanism was 10 times faster than VCG



Thank You!







LINK TO PAPER