

# Oracle Complexity of Single-Loop Switching Subgradient Methods for Non-Smooth Weakly Convex Functional Constrained Optimization

**Yankun Huang<sup>1</sup>**, Qihang Lin<sup>1</sup>

<sup>1</sup>Department of Business Analytics, University of Iowa

Supported by NSF Grant No 2147253

# Non-Smooth Weakly Convex Constrained Optimization

Problem formulation:

$$f^* \equiv \min_{\mathbf{x}} f(\mathbf{x}) \quad \text{s.t.} \quad g(\mathbf{x}) \leq 0 \quad (\text{P})$$

## Assumption 1

- $f$  and  $g$  are real-valued and  $M$ -Lipschitz continuous (but not necessarily smooth).
- $f$  and  $g$  are  $\rho$ -weakly convex (i.e.,  $f(\mathbf{x}) + \frac{\rho}{2}\|\mathbf{x}\|^2$  and  $g(\mathbf{x}) + \frac{\rho}{2}\|\mathbf{x}\|^2$  are convex).
- $\underline{f} := \inf f(\mathbf{x}) > -\infty$ .

# Near $\epsilon$ -Stationarity

Following the literature on weakly convex optimization (Davis and Drusvyatskiy, 2019, Davis and Grimmer, 2019, Ma et al., 2020, Jia and Grimmer, 2022) consider the following near  $\epsilon$ -stationarity.

## Definition

$\mathbf{x}$  is an  $\epsilon$ -stationary point if there exist  $\lambda \geq 0$ ,  $\zeta_f \in \partial f(\mathbf{x})$  and  $\zeta_g \in \partial g(\mathbf{x})$  s.t.

$$\|\zeta_f + \lambda \zeta_g\| \leq \epsilon, \quad |\lambda g(\mathbf{x})| \leq \epsilon^2, \quad g(\mathbf{x}) \leq \epsilon^2, \quad \lambda \geq 0.$$

## Definition

$\mathbf{x}$  is a nearly  $\epsilon$ -stationary point if there exists  $\hat{\mathbf{x}}$  s.t.  $\hat{\mathbf{x}}$  is an  $\epsilon$ -stationary point and  $\|\hat{\mathbf{x}} - \mathbf{x}\| \leq \epsilon$ .

# Existing Techniques

- Solving (P) means to find a nearly  $\epsilon$ -stationary point of (P).
- Existing double-loop methods (Ma et al., 2020, Boob et al., 2023, Jia and Grimmer, 2022) find a nearly  $\epsilon$ -stationary point of (P) with oracle complexity  $O(1/\epsilon^4)$  under different CQs.
- The **oracle complexity** is the total number of times for which the algorithm queries the subgradient or function value of  $f$  or  $g$ .

# Main contribution

- Study the classical **switching subgradient (SSG) method** (Polyak, 1967) and show that,
- as a single-loop first-order algorithm, SSG can also find a nearly  $\epsilon$ -stationary point of  $(P)$  with oracle complexity  $O(1/\epsilon^4)$ .
- Invent a **switching step-size rule** to accompany the switching subgradient.

# Switching Subgradient Method

---

## Algorithm 1: Switching Subgradient (SSG) method

---

```
1 Input:  $\mathbf{x}^{(0)}$ ,  $T$ , step-sizes  $\eta_t > 0$  and tolerances  $\epsilon_t \geq 0$ .
2 for  $t = 0, 1, \dots, T - 1$  do
3   if  $g(\mathbf{x}^{(t)}) \leq \epsilon_t$  then
4      $\mathbf{x}^{(t+1)} = \mathbf{x}^{(t)} - \eta_t \zeta_f^{(t)}$  for some  $\zeta_f^{(t)} \in \partial f(\mathbf{x}^{(t)})$  and,  $I = I \cup \{t\}$ .
5   else
6      $\mathbf{x}^{(t+1)} = \mathbf{x}^{(t)} - \eta_t \zeta_g^{(t)}$  for some  $\zeta_g^{(t)} \in \partial g(\mathbf{x}^{(t)})$  and,  $J = J \cup \{t\}$ .
7   end
8 end
9 Output:  $\mathbf{x}^{(\tau)}$  where  $\tau$  is sampled from  $I \cup J$  using  $\text{Prob}(\tau = t) = \eta_t / \sum_{s \in I \cup J} \eta_s$ .
```

---

# Technical Lemmas

## Assumption 2 (Uniform Slater's condition in Ma et al. (2020))

There exist  $\bar{\epsilon} > 0$ ,  $\theta > 0$  and  $\bar{\rho} > \rho$  such that Slater's condition

$$\exists \mathbf{y} \text{ s.t. } g(\mathbf{y}) + \frac{\bar{\rho}}{2} \|\mathbf{y} - \mathbf{x}\|^2 \leq -\theta$$

holds for any  $\mathbf{x}$  satisfying  $g(\mathbf{x}) \leq \bar{\epsilon}^2$ . (This is the CQ for SSG in our results.)

Denote:  $g_+(\mathbf{x}) = \max\{g(\mathbf{x}), 0\}$ ,  $\mathcal{L} = \{\mathbf{x} \mid g(\mathbf{x}) = 0\}$ ,  $\mathcal{S} = \{\mathbf{x} \mid g(\mathbf{x}) \leq 0\}$ .

## Lemma

1. Subgradient of  $g$  is bounded away from zero on  $\mathcal{L}$ :

$$\min_{\zeta_g \in \partial g(\mathbf{x})} \|\zeta_g\| \geq \nu := \sqrt{2\theta(\hat{\rho} - \rho)}, \quad \forall \mathbf{x} \in \mathcal{L} \text{ for some } \hat{\rho} \in (\rho, \bar{\rho}].$$

2. Local error bound holds:

$$(\nu/2) \cdot \text{dist}(\mathbf{x}, \mathcal{S}) \leq g_+(\mathbf{x}) \text{ if } \text{dist}(\mathbf{x}, \mathcal{S}) \leq \nu/\rho.$$

# Oracle Complexity

When  $g(\mathbf{x}^{(t)}) > \epsilon_t$ , SSG is essentially solving a sharp weakly convex unconstrained problem

$$\mathcal{S} = \arg \min_{\mathbf{x}} g_+(\mathbf{x}),$$

and thus Davis et al. (2018) suggests applying the Polyak's step-size in this case for the Q-linear convergence on  $\text{dist}(\mathbf{x}^{(t)}, \mathcal{S})$ .

## Theorem 1

Suppose  $\hat{\rho} \in (\rho, \bar{\rho}]$  and  $\epsilon \leq \bar{\epsilon}$ . Let  $\mathbf{x}^{(0)} \in \mathcal{S}$ ,  $\epsilon_t = \frac{\nu}{4} \min \{\epsilon^2/M, \nu/(4\rho)\}$  and

$$\eta_t = \begin{cases} \frac{\nu}{4M^2} \min \{\epsilon^2/M, \nu/(4\rho)\} & \text{if } g(\mathbf{x}^{(t)}) \leq \epsilon_t \\ g(\mathbf{x}^{(t)})/\|\zeta_g^{(t)}\|^2 & \text{if } g(\mathbf{x}^{(t)}) > \epsilon_t. \end{cases}$$

Then  $g(\mathbf{x}^{(t)}) \leq \epsilon^2$ ,  $\forall t \geq 0$ , and SSG finds a nearly  $\epsilon$ -stationary point of (P) if

$$T \geq \frac{8M^2 \left( f(\mathbf{x}^{(0)}) - \underline{f} + 3M^2/(2\hat{\rho}) \right)}{\hat{\rho}(1 + 2M/\nu)\nu\epsilon^2 \min \{\epsilon^2/M, \nu/(4\rho)\}} = O(1/\epsilon^4).$$

The choice of step-sizes  $\{\eta_t\}_{t \geq 0}$  in Theorem 1 shows the switching step-size rule.

- [1] Digvijay Boob, Qi Deng, and Guanghui Lan. Stochastic first-order methods for convex and nonconvex functional constrained optimization. *Mathematical Programming*, 197(1):215–279, 2023. doi:[10.1007/s10107-021-01742-y](https://doi.org/10.1007/s10107-021-01742-y).
- [2] Damek Davis and Dmitriy Drusvyatskiy. Stochastic model-based minimization of weakly convex functions. *SIAM Journal on Optimization*, 29(1):207–239, 2019. doi:[10.1137/18M1178244](https://doi.org/10.1137/18M1178244).
- [3] Damek Davis and Benjamin Grimmer. Proximally guided stochastic subgradient method for nonsmooth, nonconvex problems. *SIAM Journal on Optimization*, 29(3):1908–1930, 2019. doi:[10.1137/17M1151031](https://doi.org/10.1137/17M1151031).
- [4] Damek Davis, Dmitriy Drusvyatskiy, Kellie J. MacPhee, and Courtney Paquette. Subgradient methods for sharp weakly convex functions. *Journal of Optimization Theory and Applications*, 179(3):962–982, 2018. doi:[10.1007/s10957-018-1372-8](https://doi.org/10.1007/s10957-018-1372-8).
- [5] Zhichao Jia and Benjamin Grimmer. First-order methods for nonsmooth nonconvex functional constrained optimization with or without slater points. *arXiv preprint arXiv:2212.00927*, 2022. doi:[10.48550/arXiv.2212.00927](https://doi.org/10.48550/arXiv.2212.00927).
- [6] Runchao Ma, Qihang Lin, and Tianbao Yang. Quadratically regularized subgradient methods for weakly convex optimization with weakly convex constraints. In *International Conference on Machine Learning*, pages 6554–6564. PMLR, 2020. URL <http://proceedings.mlr.press/v119/ma20d.html>.

- [7] Boris T. Polyak. A general method for solving extremal problems. In *Doklady Akademii Nauk SSSR*, volume 174, pages 33–36, 1967. URL <https://www.mathnet.ru/eng/dan33049>.